

Contributions

- Rounding scheme for Lovász theta-function.
- Works on triangle-free perfect graphs and chordal graphs.
- No need to solve Lovász theta-function sequentially.

Maximum Stable Set Problem and Value Function Approximation.

- Maximum Stable Set Problem:

$$\max \sum_{i \in S} w_i \quad \text{s.t. } S \text{ is a stable set of } G,$$

where $G = (N, E)$ is a simple graph.

- $\mathcal{V} : 2^N \rightarrow \mathbb{R}_+$ is a *value function approximation (VFA)* if
 - it is monotone,
 - $\mathcal{V}(\emptyset) = 0$,
 - $\mathcal{V}(I) - \mathcal{V}(I \setminus (i \cup \delta(i))) \geq w_i, \forall I \subseteq N, i \in I$.

Algorithm

Algorithm 1: Retrieving a stable set from a VFA

- 1: **Input:** $G = (N, E)$, a weight function w and a VFA \mathcal{V} .
- 2: $S := \emptyset$ ▷ Start with the empty stable set
- 3: $I := \{i \in N : \mathcal{V}(i) = w_i\}$ ▷ Discard all nodes believed to be suboptimal
- 4: **while** $I \neq \emptyset$ **do**
- 5: $S := S \cup \{i\}$ for an arbitrary $i \in I$ ▷ Select i to join the stable set
- 6: $I := I \setminus (\{i\} \cup \delta_i)$ ▷ Discard i and δ_i from remaining nodes
- 7: **for** $i \in I$ **do**
- 8: **if** $\mathcal{V}(I) - \mathcal{V}(I \setminus (\{i\} \cup \delta_i)) > w_i$ **then**
- 9: $I := I \setminus \{i\}$. ▷ Discard i from remaining nodes
- 10: **Output:** Return S , a stable set of G .

LP Value Function

- **Clique LP:**

$$\max w^\top x$$

$$\text{s.t. } \sum_{i \in c} x_i \leq 1, \forall c \in C(G)$$

$$x \geq 0,$$

where $C(G)$ is the set of cliques contained in G .

$$\min \sum_{c \in C} \mu_c$$

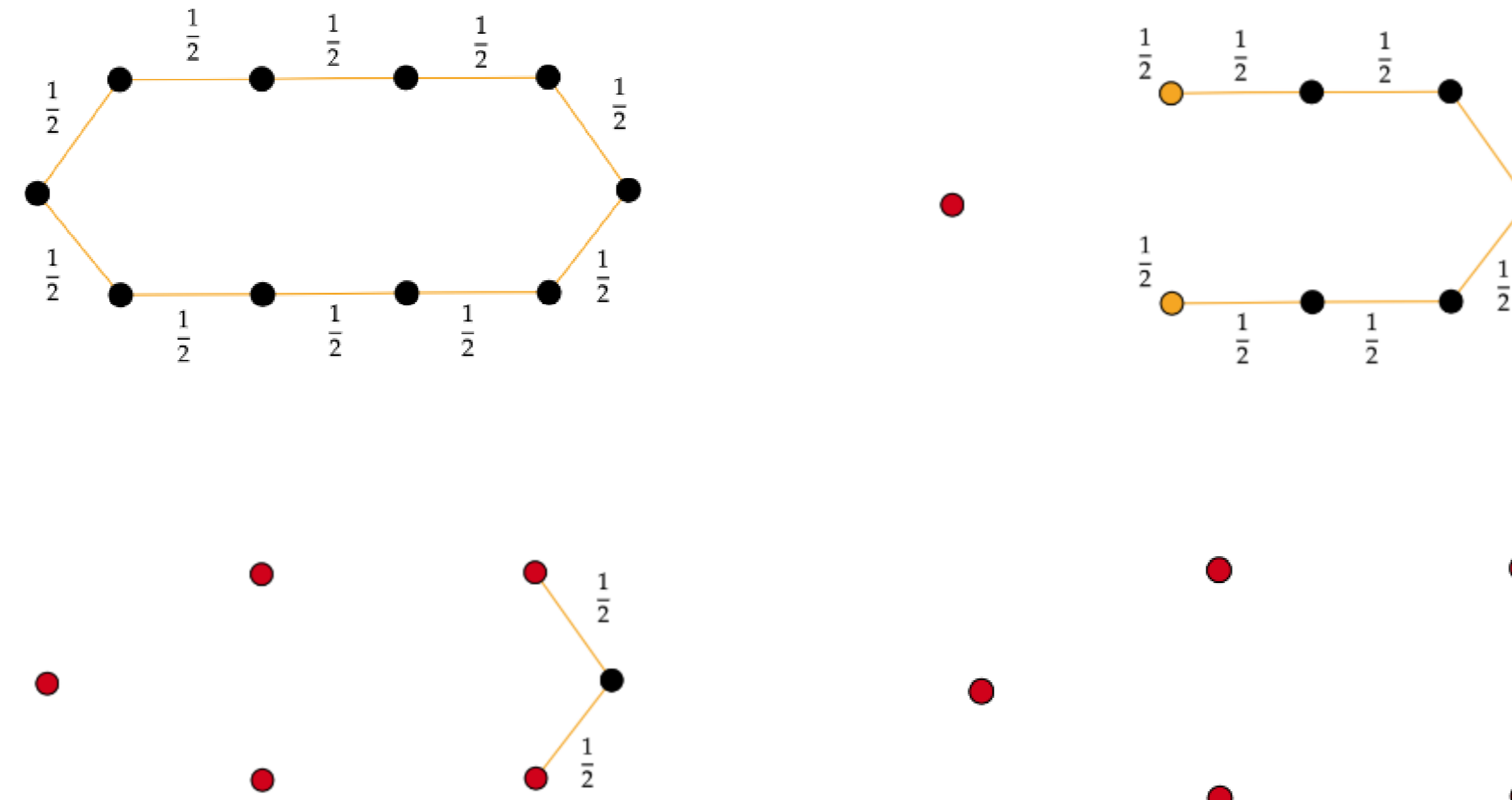
$$\text{s.t. } \sum_{c \ni i} \mu_c \geq w_i, \forall i \in N$$

$$\mu_c \geq 0.$$

- VFA based on LP: Given a pair of strictly complementary solutions (x^*, μ^*) , for $I \subseteq N$,

$$V_{LP}(I) := \sum_{c \in C(G), c \cap I \neq \emptyset} \mu_c^*$$

Illustration of Algorithm 1 with V_{LP}



SDP Value Function

- Lovász-Schrijver SDP [1]:

$$\max_{x \in \mathbb{R}^n, X \in \mathbb{S}^n} \langle \text{Diag}(w), X \rangle$$

$$\text{s.t. } X_{ii} = x_i, \forall i \in N$$

$$X_{ij} = 0, \forall ij \in E \quad (\text{SDP-P})$$

$$\begin{bmatrix} 1 & x^\top \\ x & X \end{bmatrix} \succeq 0,$$

$$\min_{t, q, Q} t$$

$$\text{s.t. } \begin{bmatrix} t & q^\top \\ q & Q \end{bmatrix} \succeq 0, \quad (\text{SDP-D})$$

$$Q_{ij} := -2q_i - w_i, \forall i = j$$

$$Q_{ij} = 0, \forall ij \notin E.$$

- Given a relative-interior optimal solution (t, q, Q) for the dual problem (SDP-D),

$$V_{SDP}(I) := q_I^\top Q_I^\dagger q_I,$$

where Q^\dagger denotes the pseudo-inverse of Q .

- Algorithm 1 evaluates the function V_{SDP} at most $O(n)$ times, and produces a stable set $S \subseteq N$ for an arbitrary graph G . Notice that q, Q are fixed; **only one SDP is solved**, unlike many existing algorithms [2]–[4].

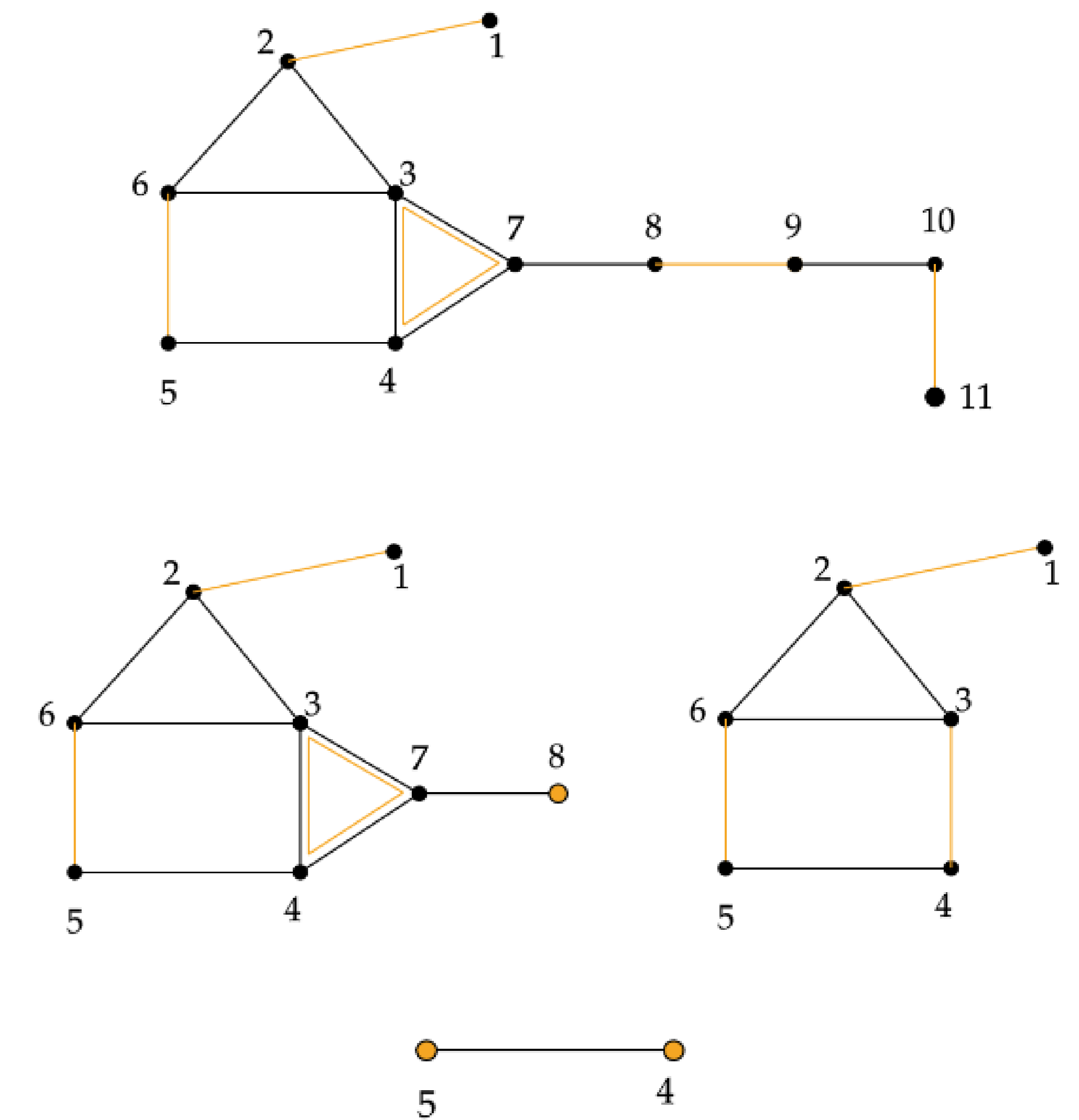
Theorem 1

Algorithm 1 with V_{LP} outputs a maximum stable set for all triangle-free perfect graphs and chordal graphs.

Theorem 2

Algorithm 1 with V_{SDP} is at least as good as the one with V_{LP} . In particular, it outputs a maximum stable set for triangle-free perfect graphs and chordal graphs.

Failing Example



References

- [1] L. Lovász and A. Schrijver, "Cones of matrices and set-functions and 0–1 optimization," *SIAM Journal on Optimization*, vol. 1, no. 2, pp. 166–190, 1991.
- [2] F. Alizadeh, "A sublinear-time randomized parallel algorithm for the maximum clique problem in perfect graphs," in *Proceedings of the Second Annual ACM-SIAM Symposium on Discrete Algorithms*, ser. SODA '91, San Francisco, California, USA: Society for Industrial and Applied Mathematics, 1991, pp. 188–194.
- [3] E. A. Yildirim and X. Fan-Orzechowski, "On Extracting Maximum Stable Sets in Perfect Graphs Using Lovász's Theta Function," *Computational Optimization and Applications*, vol. 33, no. 2, pp. 229–247, 2006.
- [4] M. Grötschel, L. Lovász, and A. Schrijver, "Polynomial algorithms for perfect graphs," in *Topics on Perfect Graphs*, ser. North-Holland Mathematics Studies, vol. 88, North-Holland, 1984, pp. 325–356.